

On Functional Logic Programming and its Application to Testing

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Software Correctness

Astronauts died because of programming errors:

- In *2001: A Space Odyssey* (by Arthur C. Clarke) computer killed astronauts because of **programming contradiction**
- In 1996, unmanned *Ariane 5* rocket exploded on its first flight because of **error in software design**

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How to avoid errors in software:

- prove program correct regarding specification
- test program properties with example input

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How to avoid errors in software:

- prove program correct regarding specification
- test program properties with example input
- use declarative programming language

Even declarative programs can contain bugs!

Imperative Programming

```
// reverse an array beginning at position 0
pos = 0;

// loop through whole array
while (pos < array.size) {

    // swap two elements
    elem = array[pos];
    array[pos] = array[array.size - pos];
    array[array.size - pos] = elem;

    // advance to next position
    pos = pos + 1;
}
```

Imperative Programming

```
// reverse an array beginning at position 0
pos = 0;

// loop through whole array
while (pos < array.size) {

    // swap two elements                off-by-one error!
    elem = array[pos];
    array[pos] = array[array.size - pos];
    array[array.size - pos] = elem;

    // advance to next position
    pos = pos + 1;
}
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}
```

Imperative Programming

```
pos    = 0  
array = [ 1 , 2 , 3 , 4 ]
```

Imperative Programming

```
pos    = 0  
array = [ 1 , 2 , 3 , 4 ]
```

```
pos    = 1  
array = [ 4 , 2 , 3 , 1 ]
```

Imperative Programming

```
pos    = 0  
array  = [ 1 , 2 , 3 , 4 ]
```

```
pos    = 1  
array  = [ 4 , 2 , 3 , 1 ]
```

```
pos    = 2  
array  = [ 4 , 3 , 2 , 1 ]
```

Imperative Programming

```
pos    = 0  
array = [ 1 , 2 , 3 , 4 ]
```

```
pos    = 1  
array = [ 4 , 2 , 3 , 1 ]
```

```
pos    = 2  
array = [ 4 , 3 , 2 , 1 ]
```

```
pos    = 3  
array = [ 4 , 2 , 3 , 1 ]
```

Imperative Programming

```
pos = 0  
array = [ 1 , 2 , 3 , 4 ]
```

```
pos = 1  
array = [ 4 , 2 , 3 , 1 ]
```

```
pos = 2  
array = [ 4 , 3 , 2 , 1 ] ← stop here!
```

```
pos = 3  
array = [ 4 , 2 , 3 , 1 ]
```

```
pos = 4  
array = [ 1 , 2 , 3 , 4 ]
```

Imperative Programming

```
// reverse an array beginning at position 0
pos = 0;

// loop through half array
while (pos < array.size / 2) {

    // swap two elements
    elem = array[pos];
    array[pos] = array[array.size - pos - 1];
    array[array.size - pos - 1] = elem;

    // advance to next position
    pos = pos + 1;
}
```

Imperative Programming

- **program is sequence of statements**
- **built-in control structures (loops, conditional branches)**
- **values of variables change (side effects)**
- **results depend on evaluation order**
- **easy to make mistakes (index manipulations)**

Functional Programming

Language: Haskell

reverse [] = []

reverse (x : xs) = reverse xs ++ (x : [])

Functional Programming

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[] ++ ys = ys
(x : xs) ++ ys = x : (xs ++ ys)

Functional Programming

reverse [] = []

reverse (x : xs) = **reverse** xs ++ (x : [])

reverse (1 : 2 : 3 : 4 : [])

= **reverse** (2 : 3 : 4 : [])

++ (1 : [])

Functional Programming

reverse [] = []

reverse (x : xs) = **reverse** xs ++ (x : [])

reverse (1 : 2 : 3 : 4 : [])

= **reverse** (2 : 3 : 4 : []) ++ (1 : [])

= **reverse** (3 : 4 : []) ++ (2 : []) ++ (1 : [])

= **reverse** (4 : []) ++ (3 : []) ++ (2 : []) ++ (1 : [])

= **reverse** [] ++ (4 : []) ++ (3 : []) ++ (2 : []) ++ (1 : [])

= 4 : 3 : 2 : 1 : []

Functional Programming

- **program is set of equations between expressions**
- **recursion instead of built-in control structures**
- **values of variables do not change**
- **results independent of evaluation order**
- **constructors create, patterns eliminate data**

Functional Logic Programming

Language: Curry

last list

| **list** \equiv **xs** ++ [**x**] – guard for the equation defining ‘last’
= **x**

where **x**, **xs** free – free variables for unknown values

Functional Logic Programming

Language: Curry

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| **list** \equiv **xs** ++ [**x**] – guard for the equation defining ‘last’
= **x**

where **x**, **xs** free – free variables for unknown values

if **list** = [1, 2, 3, 4] guard is satisfied for **xs** = [1, 2, 3] and **x** = 4:

$$[1, 2, 3, 4] \equiv [1, 2, 3] ++ [4]$$

$$\rightsquigarrow \mathbf{last} [1, 2, 3, 4] = 4$$

Functional Logic Programming

insert x xs | xs \equiv ys ++ zs
= ys ++ [x] ++ zs
where **ys, zs** free

Functional Logic Programming

insert x xs | $xs \equiv ys ++ zs$
 $= ys ++ [x] ++ zs$
where ys, zs free

insert 1 [2, 3, 4]
 $=$ **insert** 1 ([2] ++ [3, 4])
 $=$ [2] ++ [1] ++ [3, 4]
 $=$ [2, 1, 3, 4]

Functional Logic Programming

insert x xs | **xs** \equiv **ys** ++ **zs**
= **ys** ++ [**x**] ++ **zs**
where **ys, zs** free

insert 1 [2, 3, 4]
= **insert 1 ([2] ++ [3, 4])**
= [2] ++ [1] ++ [3, 4]
= [2, 1, 3, 4]

insert 1 [2, 3, 4]
= **insert 1 ([2, 3] ++ [4])**
= [2, 3] ++ [1] ++ [4]
= [2, 3, 1, 4]

Functional Logic Programming

- **free variables for unknown values**
- **built-in search for solving equational guards**
- **nondeterministic results**

Narrowing

```
cyi> reverse xs where xs free
```

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{xs = []} []
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```
More solutions? [Y(es)/n(o)/a(ll)] yes
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{xs = [_a]} [_a]
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```
{xs = [_a]} [_a]
```

```
More solutions? [Y(es)/n(o)/a(11)] yes
```

```
{xs = [_b,_c]} [_c,_b]
```

```
More solutions? [Y(es)/n(o)/a(11)] yes
```

Narrowing

cyi> reverse xs where xs free

{xs = []} []

More solutions? [Y(es)/n(o)/a(11)] yes

{xs = [_a]} [_a]

More solutions? [Y(es)/n(o)/a(11)] yes

{xs = [_b,_c]} [_c,_b]

More solutions? [Y(es)/n(o)/a(11)] yes

{xs = [_d,_e,_f]} [_f,_e,_d]

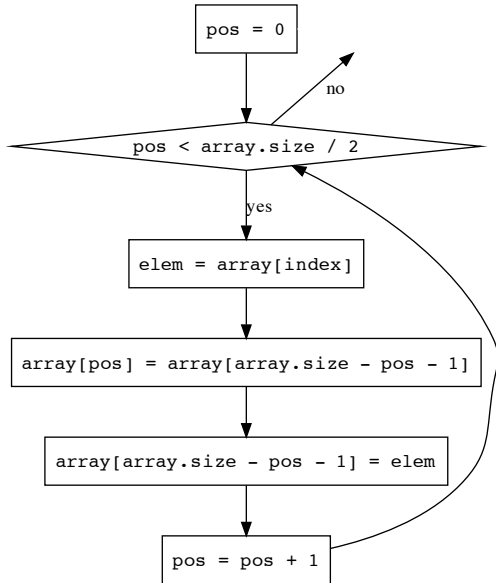
More solutions? [Y(es)/n(o)/a(11)] no

Narrowing

- **generates test cases for free** (F. and Kuchen, PPDP 2007)
- **often infinitely many**
- **which tests are *redundant*?**

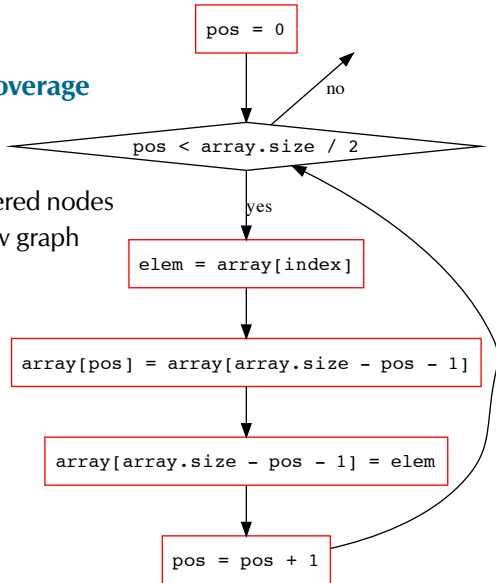
Code Coverage

- **groups tests into equivalence classes**
- **maps all possible program behaviours to finite set of classes**
- **approximation of program behaviour**
 - too fine: similar tests considered different
 - too coarse: interesting tests considered redundant
- **different notions, depending on paradigm**



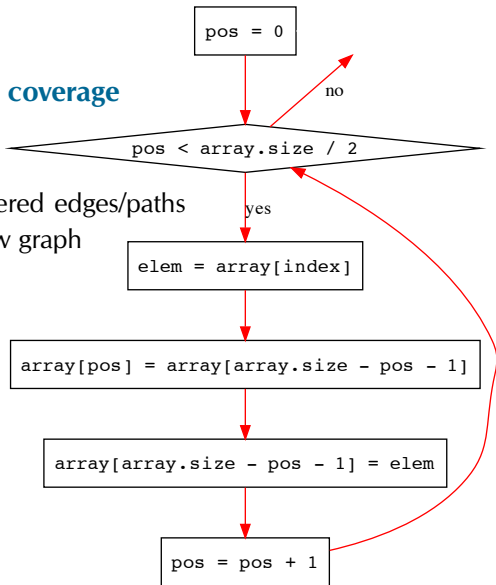
statement coverage

monitors covered nodes
in control-flow graph



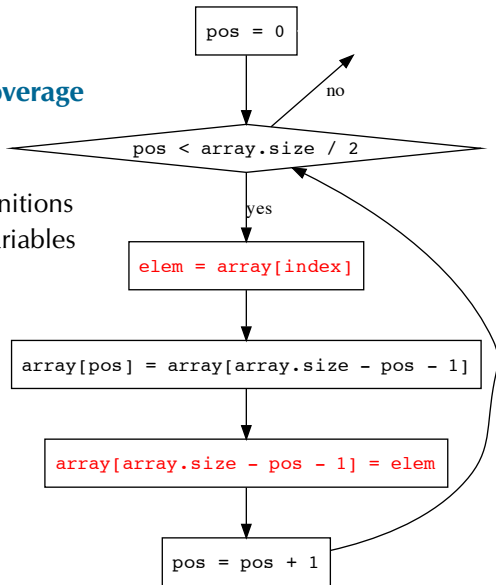
control-flow coverage

monitors covered edges/paths
in control-flow graph



data-flow coverage

monitors definitions
and uses of variables



Declarative Code Coverage

criteria from imperative programming cannot be transferred easily

- no sequence of statements
- no modifiable variables

new notions of code coverage for declarative programs

- **expression coverage**
- **control flow** (F. and Kuchen, PPDP 2007)
- **data flow** (F. and Kuchen, ICFP 2008)

Expression Coverage

which expressions need to be evaluated?

test xs ys = reverse (xs ++ ys) \equiv reverse ys ++ reverse xs

reverse [] = []

reverse (x : xs) = [x] ++ reverse xs

[] ++ ys = ys

(x : xs) ++ ys = x : (xs ++ ys)

test [1] [] demands evaluation of all expressions in program

Expression Coverage

which expressions need to be evaluated?

test xs ys = reverse (xs ++ ys) \equiv reverse ys ++ reverse xs

reverse [] = []

reverse (x : xs) = [x] ++ reverse xs

[] ++ ys = ys

(x : xs) ++ ys = x : (xs ++ ys)

test [1] [2] would have revealed a **bug** \rightsquigarrow criterion too coarse

Control-Flow Coverage

used equations of (buggy) **reverse** for evaluation of **test** [1] []:

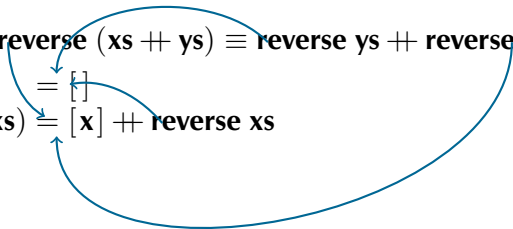
test xs ys = **reverse** (xs ++ ys) \equiv **reverse** ys ++ **reverse** xs

reverse [] = []

reverse (x : xs) = [x] ++ **reverse** xs

function call

equation



Control-Flow Coverage

used equations of (buggy) **reverse** for evaluation of **test** [1] [2]:

test *xs* *ys* = **reverse** (*xs* ++ *ys*) \equiv **reverse** *ys* ++ **reverse** *xs*

reverse []

= []

reverse (*x* : *xs*) = [*x*] ++ **reverse** *xs*

different arrows

function call

equation



Control-Flow Coverage

full control-flow coverage for (buggy) **reverse**:

test [] [] = **True**

test [1] [2] = **False**

Control-Flow Coverage

full control-flow coverage for (buggy) **reverse**:

test [] [] = **True**

test [1] [2] = **False**

test [] [] = **True**

test [1] [] = **True**

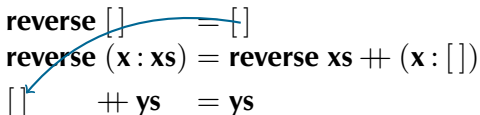
test [] [1, 2] = **True**

errors may remain undetected despite full coverage!

Data-Flow Coverage

data flow for (correct) evaluation of **reverse** [1]:

reverse [] = []
reverse (x : xs) = **reverse** xs ++ (x : [])
[] ++ ys = ys
(x : xs) ++ ys = x : (xs ++ ys)



A blue curved arrow originates from the first `[]` in the second equation, `reverse (x : xs) = reverse xs ++ (x : [])`, and points to the `[]` in the third equation, `[] ++ ys = ys`.

constructor call

pattern



A blue curved arrow originates from the text "constructor call" and points to the text "pattern".

Data-Flow Coverage

data flow for (correct) evaluation of **reverse** [1, 2]:

reverse [] = []
reverse (x : xs) = **reverse** xs ++ (x : [])
[] ++ ys = ys
(x : xs) ++ ys = x : (xs ++ ys)

The diagram consists of four lines of code. Blue arrows indicate data flow from patterns to constructors. One arrow points from the pattern '[]' in the second line to the constructor '[]' in the first line. Another arrow points from the pattern '[]' in the third line to the constructor '[]' in the second line. A third arrow points from the pattern 'xs' in the fourth line to the constructor 'xs' in the second line. A fourth arrow points from the pattern 'xs' in the fourth line to the constructor 'xs' in the third line.

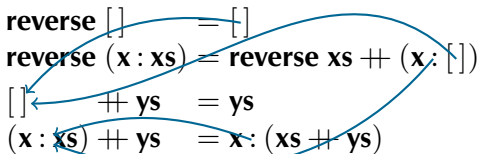
constructor call

pattern

Data-Flow Coverage

data flow for (correct) evaluation of **reverse** [1, 2, 3]:

reverse [] = []
reverse (x : xs) = **reverse** xs ++ (x : [])
[] ++ ys = ys
(x : ~~xs~~) ++ ys = x : (xs ++ ys)



constructor call

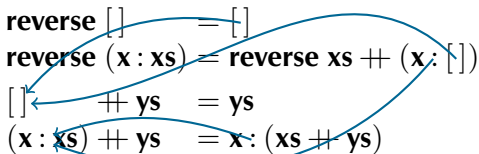
pattern



Data-Flow Coverage

data flow for (correct) evaluation of **reverse** [1, 2, 3, 4]:

reverse [] = []
reverse (x : xs) = **reverse** xs ++ (x : [])
[] ++ ys = ys
(x : ~~xs~~) ++ ys = x : (xs ++ ys)



no new data flow for length 4

constructor call

pattern



Experiments

- introduce errors in algorithmically challenging programs
 - AVL tree (insert and delete functions)
 - Heapsort with purely functional heap
 - Strassen's algorithm for matrix multiplication
 - Dijkstra's shortest path algorithm
 - Kruskal's minimum spanning tree algorithm
 - Matrix Chain Multiplication (dynamic programming)
- generate tests automatically
- eliminate redundant tests regarding different coverage criteria
- check whether failing tests remain

Experiments

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 - AVL tree (insert and delete functions)
 - Heapsort with purely functional heap
 - Strassen's algorithm for matrix multiplication
 - Dijkstra's shortest path algorithm
 - Kruskal's minimum spanning tree algorithm ← **textbook error**
 - Matrix Chain Multiplication (dynamic programming)
- generate tests automatically
- eliminate redundant tests regarding different coverage criteria
- check whether failing tests remain

Experiments

	Control Flow	Data Flow	Control + Data Flow
AVL insert	9	12	20
AVL delete	9	23	27
Heapsort	6	7	16
Strassen	4	13	18
Dijkstra	6	9	10
Kruskal	4	9	10
Matrix CM	2	7	4

Table: Sizes of reduced sets of tests for different coverage criteria

Declarative Code Coverage

used to classify tests for algorithmically challenging program units

control-flow coverage

- which equations are used by function calls?
- distinguishes different calls that use an equation

data-flow coverage

- which constructors are matched in patterns?
- distinguishes different values matched in an equation

combination is more thorough than one criterion alone

Thesis Contents

- **Declarative Programming**
 - **Functional programming**
 - **Functional logic programming**
- Generating Tests
 - Black-box testing
 - Glass-box testing
- **Code coverage**
 - **Control flow**
 - **Data flow**
 - Monitoring code coverage
 - **Experimental evaluation**
- Explicit nondeterminism
 - Nondeterminism monads
 - Combining laziness with nondeterminism

F. and Kuchen **PPDP 2007. Systematic generation of glass-box test cases for functional logic programs**

Proceedings of the 9th ACM SIGPLAN International Symposium on Principles and Practice of Declarative Programming. ACM Press.

Christiansen and F. **FLOPS 2008. EasyCheck – test data for free**

Proceedings of the 9th International Symposium on Functional and Logic Programming. Springer LNCS 4989.

F. and Kuchen **ICFP 2008. Data-flow testing of declarative programs**

Proceedings of the 13th ACM SIGPLAN International Conference on Functional Programming. ACM Press.

F. **ATPS 2009. Reinventing Haskell backtracking**

Proceedings der GI-Jahrestagung Informatik 2009. GI LNI.

F., Kiselyov, and Shan **ICFP 2009.**

Purely functional lazy non-deterministic programming

Proceedings of the 14th ACM SIGPLAN International Conference on Functional Programming. ACM Press.