

Declarative Programming of User Interfaces^{*}

(Extended Abstract)

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Abstract. This paper proposes a declarative description of user interfaces that abstracts from low-level implementation details. In particular, the user interfaces specified in our framework are executable as graphical user interfaces on single local host computers as well as web user interfaces via web standard browsers. Thus, our approach combines the advantages of existing user interface technologies in a flexible way without demands on the programmer's side. We sketch an implementation of this concept in the declarative multi-paradigm programming language Curry and show how the integrated functional and logic features of Curry are exploited to enable a high-level implementation of this concept.

1 Motivation

The implementation of a good user interface for application programs is a necessary but often non-trivial and tedious task. In order to support programmers in the implementation of user interfaces, one can find specific libraries that reflect different approaches to the construction of user interfaces. From a user's perspective, there are two kinds of user interfaces (UIs) that are currently the most important ones on conventional desktop computers:

Graphical User Interfaces (GUIs): These are user interfaces that followed the early textual user interfaces on single host computers. GUIs enabled non-expert users to easily interact with application programs. They provide a good reaction time (since they run on the local host) and are relatively easy to install as any other program, i.e., usually they are distributed with the executable of the application program. On the negative side, application programs with GUIs are not so easy to install for many users (since one has to install them on all local hosts) and are difficult to maintain during their life time (since updates must be performed on all existing installations).

Web User Interfaces (WUIs): These are user interfaces that became popular with the world-wide web and its opportunities for user interaction via dynamic web pages. In this case, the application runs on a web server and the user interacts with the application via a standard web browser. Thus,

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applications with WUIs are relatively easy to install for many users since every single user needs only a web browser on his local host (which is usually already installed). Moreover, such applications are easy to maintain since one only has to update the central installation on the web server. On the negative side, WUIs have a moderate reaction time (since the web server is contacted for every state-changing interaction) and a complete application is more difficult to install on a single host (since one has to install and configure a web server).

A few years ago, there was also another important difference between GUIs and WUIs: the model of interaction. In application with GUIs, the user could immediately change the content of many widgets by mouse events, whereas with WUIs, each page containing user input has to be sent to the web server which sends back a new web page with some modified content. However, this disadvantage of WUIs has been decreased or omitted by the development of the Ajax framework that supports an interaction with a web server without submitting and receiving complete new pages from the web server [6].

From these considerations, it is reasonable to combine the advantages of both kinds of user interfaces in a single framework so that the programmer has no additional burden to select between GUIs or WUIs (or both) for his application. This paper presents a concrete proposal of such a concept and its implementation in the declarative multi-paradigm language Curry.

Adobe AIR¹ also enables the use of the same base code to create applications that run in a web browser as well as on a desktop. In contrast to our approach, Adobe AIR is not based on standard features of web browsers but requires specific software to be installed on the client's side. Another related work is the Google Web Toolkit (GWT²). GWT is a framework to implement dynamic web pages for Java programs similarly to GUI programming in order to create highly interactive web applications with reasonable efforts. In contrast to our proposal, GWT does not support the use of the same program to generate both GUI and WUI applications. Moreover, our framework is based on declarative programming concepts so that we exploit the abstract programming techniques provided by declarative languages to support a high-level programming of user interfaces. In our framework, one specifies the structure of user interfaces from which the concrete implementations (GUIs or WUIs) are automatically inferred.

In the following section, we review the main features of functional logic programming as required in this paper. Section 3 describes the concepts of our framework followed by a few examples shown in Section 4. Implementation issues are sketched in Sections 5 and 6 before we conclude in Section 7.

¹ <http://www.adobe.com/devnet/air/>

² <http://code.google.com/webtoolkit/>

2 Functional Logic Programming and Curry

We assume familiarity with functional logic programming (see [13] for a recent survey) and Curry [7, 17]. Therefore, we review only the basic concepts relevant for this paper.

Functional logic languages integrate the most important features of functional and logic languages to provide a variety of programming concepts to the programmer. Modern languages of this kind [7, 17, 18] combine the concepts of demand-driven evaluation and higher-order functions from functional programming with logic programming features like computing with partial information (logic variables), unification, and nondeterministic search for solutions. This combination, supported by optimal evaluation strategies [1] and new design patterns [2], leads to better abstractions in application programs, e.g., as shown for programming with databases [3, 5, 10] or web programming [9, 12, 14]. The declarative multi-paradigm Curry [7, 17] is a functional logic language extended by concurrent programming concepts and has been used in various applications. In the following, we review the elements of Curry that are relevant to understand the contents of this paper. Further features (e.g., constraints, search strategies, concurrency, I/O concept, modules), more details about Curry's computation model and a complete description of the language can be found in [7, 17].

From a syntactic point of view, a Curry program is a functional program extended by the possible inclusion of free (logic) variables in conditions and right-hand sides of defining rules. Curry has a Haskell-like syntax [21], i.e., (type) variables and function names usually start with lowercase letters and the names of type and data constructors start with an uppercase letter. The application of f to e is denoted by juxtaposition (“ $f e$ ”). A Curry *program* consists of the definition of functions and data types on which the functions operate. Functions are first-class citizens and evaluated lazily. To provide the full power of logic programming, functions can be called with partially instantiated arguments and defined by conditional equations with constraints in the conditions. Function calls with free variables are evaluated by a possibly nondeterministic instantiation of demanded arguments (i.e., arguments whose values are necessary to decide the applicability of a rule) to the required values in order to apply a rule.

In general, functions are defined by (*conditional*) *rules* of the form “ $f t_1 \dots t_n \mid c = e$ ” with f being a function, t_1, \dots, t_n *patterns* (i.e., expressions without defined functions) without multiple occurrences of a variable, the (optional) *condition* c is a constraint (e.g., a conjunction of equations), and e is a well-formed *expression* which may also contain function calls, lambda abstractions etc. A conditional rule can be applied if its left-hand side matches the current call and its condition is satisfiable.

The following Curry program defines the data types of Boolean values, possible values, and polymorphic lists, and functions to compute the concatenation of lists and the last element of a list:

```
infixr 5 ++
data Bool    = True    | False
```

```

data Maybe a = Nothing | Just a
data List a = [] | a : List a

(++) :: [a] -> [a] -> [a]
[]      ++ ys = ys
(x:xs) ++ ys = x : (xs ++ ys)

last :: [a] -> a
last xs | ys ++ [x] ::= xs = x   where x,ys free

```

`[]` (empty list) and `:` (non-empty list) are the constructors for polymorphic lists (`a` is a type variable ranging over all types and the type “`List a`” is written as `[a]` for conformity with Haskell). The *infix operator declaration* “`infixr 5 ++`” declares the symbol “`++`” as a right-associative infix operator with precedence 5 so that we can write function applications of this symbol with the convenient infix notation. The (optional) type declaration (“`::`”) of the function “`++`” specifies that “`++`” takes two lists as input and produces an output list, where all list elements are of the same (unspecified) type.³

As one can see in this example, logic programming is supported by admitting function calls with free variables (see “`ys ++ [x]`” above) and constraints in the condition of a defining rule. For instance, the equation “`ys ++ [x] ::= xs`” is solved by instantiating the first argument `ys` to the list `xs` without the last argument, i.e., the only solution to this equation satisfies that `x` is the last element of `xs`. In order to support some consistency checks, *extra variables*, i.e., variables of a rule not occurring in a pattern of the left-hand side, must be declared by “`where . . . free`” (see the rule defining `last`). In general, a *constraint* is any expression of the built-in type `Success`. For instance, the trivial constraint `success` is an expression of type `Success` that denotes the always satisfiable constraint. “`c1 & c2`” denotes the *concurrent conjunction* of the constraints `c1` and `c2`, i.e., this expression is evaluated by proving both argument constraints concurrently. An *equational constraint* `e1 ::= e2` is satisfiable if both sides `e1` and `e2` are reducible to unifiable constructor terms. Specific Curry systems also support more powerful constraint structures, like arithmetic constraints on real numbers or finite domain constraints (e.g., PAKCS [15]).

The operational semantics of Curry, described in detail in [7, 17], is based on an optimal evaluation strategy [1] which is a conservative extension of lazy functional programming and (concurrent) logic programming. Curry also offers standard features of functional languages, like higher-order functions, modules, or monadic I/O (which is identical to Haskell’s I/O concept [24]). Thus, “`I0 a`” denotes the type of an I/O action that returns values of type “`a`”. For instance, the predefined I/O action `getChar` has the type “`I0 Char`”, i.e., it returns the next character from the keyboard when it is applied. Similarly, the predefined I/O action `readFile` has the type “`String -> I0 String`”, i.e., it takes a string (the name of a file) and returns the contents of the file when it is applied.

³ Curry uses curried function types where $\alpha \rightarrow \beta$ denotes the type of all functions mapping elements of type α into elements of type β .

3 Specifying User Interfaces

In this section we describe our framework for the declarative programming of user interfaces that can be executed either on a local host as a GUI (e.g., by the use of Tcl/Tk [20]) or as a WUI on a web server that is accessed by a standard web browser.

In order to develop appropriate abstractions for high-level UI programming, one has to analyze the essential components of these programming tasks. Based on earlier work on programming GUIs and WUIs with functional logic languages [8, 9, 12], one can distinguish the following ingredients of UI programming:

Structure: Each UI has a specific hierarchical structure which typically consists of basic elements (also called *widgets*), like text input fields or selection boxes, and composed elements, like rows or columns of widgets. Thus, UIs have a tree-like structure which can be easily specified by an algebraic data type in a declarative language.

Functionality: If the user interacts with UI elements by mouse or keyboard clicks, these UI elements emit some events on which the application program should react. A convenient way to connect the application program to such events is the concept of *event handlers*, i.e., functions that are associated to events of some widget and that are called whenever such an event occurs. Usually, the event handlers use the functionality of the application program to compute some data that is shown in the widgets of the UI. Thus, event handlers are associated to some widgets but need to refer to other widgets independently of the structural hierarchy. This means that UIs have not only a hierarchical (layout) structure but also a logical (graph-like) structure that connects the event handlers with various widgets of the UI structure. In previous works on GUI and WUI programming [8, 9] it has been shown that free (logic) variables are an appropriate feature to describe this logical structure and to avoid many problems that occur if fixed strings are used as references to UI elements as in traditional GUI programming (e.g., [20, 23]) or WUI programming (e.g., [4, 19]).

Layout: In order to support a visually appealing appearance of a UI, it should be possible to influence the standard layout of a UI. Whereas in older approaches layout and structural information are often mixed (e.g., as in Tcl/Tk or older versions of HTML, and similarly in previous approaches to declarative GUI/WUI programming [8, 9]), it has been realized that these issues should be distinguished in order to obtain clearer and reusable implementations. For instance, current versions of HTML recommend the use of cascading style sheets (CSS) to separate structure from layout.

The distinction between structure, functionality, and layout and their appropriate modelling in a declarative programming language are the key ingredients to our framework for UI programming. Although parts of these ideas can be found in our previous works [8, 9, 12], our current novel approach abstracts more from the underlying technology (Tcl/Tk, HTML/CGI) so that it enables a common

method to specify user interfaces. In the following, we propose a concrete description of the structure, functionality, and layout of UIs in the language Curry by presenting appropriate data types and operations on them. In principle, one can transfer these ideas also to other declarative languages (where some restrictions might be necessary). However, we will see that the combined functional and logic programming features of Curry are exploited for our high-level and application-oriented description of UIs.

As discussed above, UIs have a hierarchical structure that can be appropriately described by the following data type:

```
data Widget r e = Widget WidgetKind
                  (Maybe String)
                  (Maybe (Ref r))
                  [Handler e]
                  [StyleClass]
                  [Widget r e]
```

In order to avoid unnecessary restrictions, the definition of a widget is quite general. In principle, one could also enumerate all kinds of widgets and distinguish between widgets having no structure (basic widgets) and widgets with structure (e.g., rows, columns). For the sake of generality, we have chosen one widget constructor where the concrete kind of widget is given as the first component (of type `WidgetKind`). The last two components are a list of layout elements (see below) and the widgets contained in this widget, respectively. The second component contains the possible contents of the widget (e.g., the entry string of a text input field, `Nothing` for widget combinators like row or column), the third component a possible reference to a widget used by other event handlers, and the fourth component a list of handlers for the various events that can occur in this widget. The type parameter `r` varies over the possible reference types used in concrete instantiations (e.g., `Tcl/Tk` references in GUIs, `CGI` references in HTML), and the type parameter `e` varies over the possible environments passed to event handlers in concrete instantiations. These type parameters are only necessary to provide different instantiations of general UIs and need not to be considered by the programmer. Concrete examples for widgets are shown below after we have discussed the other data types used in widgets.

As discussed above, event handlers need to refer to other widgets independently of the widget hierarchy. Therefore, a widget can be equipped with an identity used as a reference by event handlers. Many approaches to user interfaces programming, like `Tcl/Tk` or `HTML/CGI`, use string constants as identifiers. Since the use of concrete strings is error prone (e.g., if strings do not exist due to typos in the program), we adapt the idea of previous works on declarative GUI and WUI programming [8, 9] to make the type of widget references abstract so that one cannot construct “wrong” identifiers but use free variables (whose declarations are checked at compile time) for this purpose. Thus, our UI library contains the type declaration

```
data Ref r = Ref r
```

where only the type name but not its data constructor is exported, i.e., `Ref` is an abstract type. Since no constructor of this data type is available to the user of the UI library, the only reasonable way to use values of type `Ref` is with a free variable (see below for a concrete example).

In general, event handlers are used for two main purposes. Either they should perform some calculations and show their results in some specific widgets of the UI, i.e., they influence the state of the UI, or they should change the state of the underlying application program, e.g., the execution of an event handler might change some application data that is stored in a file or database. In order to support the latter functionality, the result type of an event handler is always “`IO ()`”, i.e., an event handler might have a side effect on the external world. Since there are also I/O actions to influence the state of the UI (see below), this result type of event handlers ensures that event handlers can influence the state of the UI as well as the state of the application program.

Furthermore, the calculations and actions performed by event handlers usually depend on the user inputs stored in the widgets of the interface, i.e., these input values must be passed as parameters to the event handlers. This can be adequately modelled by an *environment* parameter that is conceptually a mapping from widget references to the string values stored in the widgets. Altogether, the UI library contains the type declarations

```
data Command e = Cmd (e -> IO ())
data Handler e = Handler Event (Command e)
```

where `Event` is the type of possible events issued by user interfaces:

```
data Event = DefaultEvent | FocusIn | FocusOut
           | MouseButton1 | MouseButton2 | MouseButton3
           | KeyPress | Return | Change | DoubleClick
```

Therefore, each element in the list of event handlers of a widget specifies a command (an I/O action depending on the value of some environment `e`) that is executed whenever the associated event occurs.

The type `WidgetKind` specifies the different kinds of widgets supported by our library. Some constructors of this type are

```
data WidgetKind = Col | Row | Label | Button | Entry
                | TextEdit Int Int | ...
```

The constructors `Col` and `Row` specify combinations of widgets as columns and rows, respectively. `Label` is a widget containing a string not modifiable by the user, `Button` is a simple button, `Entry` is an entry field for a line of text, and `TextEdit` is a widget to edit larger text areas (the parameters are the height and width of the edit area).

Since it would be tedious to define all widgets of a user interface by using the constructor `Widget` only, the library contains a number of useful abbreviations, like

```

col ws      = Widget Col  Nothing  Nothing  [] [] ws
row ws      = Widget Row  Nothing  Nothing  [] [] ws
label str   = Widget Label (Just str) Nothing  [] [] []
entry ref str = Widget Entry (Just str) (Just ref) [] [] []
button cmd label =
  Widget Button (Just label) Nothing
    [Handler DefaultEvent (Cmd cmd)] [] []

```

For instance, a simple UI showing the text “Hello World!” and a button to exit the UI can be specified as follows:

```

col [label "Hello World!",
     button exitUI "Stop"]

```

`exitUI` is a predefined event handler to terminate the UI. The environment passed to event handlers can be accessed and modified by the predefined I/O actions `getValue` and `setValue` that take a widget reference as their first argument. Thus, “`getValue r e`” returns the value of the widget referenced by `r` w.r.t. environment `e`, and “`setValue r v e`” updates the value of the widget referenced by `r` so that it becomes visible to the user.

In order to influence the layout of UIs, widgets can take a list of style parameters of type `StyleClass`. This type contains options to align the widget or the text contained in it, set the font and color of the widget’s text, set the background color, and so on. The styles of a widget can be dynamically changed by predefined operations like `setStyles`, `addStyles`, etc.

4 Examples

In order to demonstrate the concrete application of our concept, we show a few programming examples in this section. As a first example, consider a simple counter UI shown in Fig. 1. Using our library, its structure and functionality is specified as follows:

```

counterUI = col [label "A simple counter:",
                 entry val "0",
                 row [button inc    "Increment",
                     button reset  "Reset",
                     button exitUI "Stop" ]]
where val free
  reset env = setValue val "0" env
  inc env = do v <- getValue val env
             setValue val (show (readInt v + 1)) env

```

The free variable `val` denotes the reference to the entry field containing the string representation of the counter’s value. It is used by the event handler `reset` to set

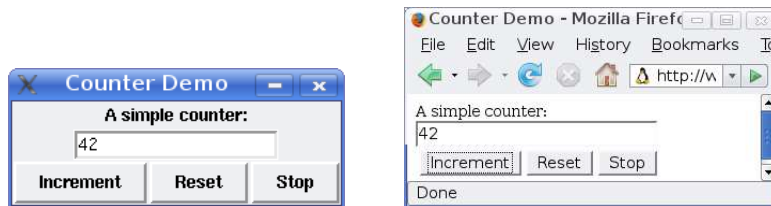


Fig. 1. A simple counter UI executed as a GUI (left) and as a WUI (right)

the value of this entry widget to "0". The event handler `inc` reads the current value of this widget before replacing it by its incremented value.

The UI specification can be executed by the predefined I/O action `runUI` that takes a string (usually shown as the label of the window containing the UI) and a UI specification as parameters. For instance, the counter UI shown above is executed by evaluating the main expression

```
runUI "Counter Demo" counterUI
```

Many interactive applications contain a state which is shown and modified by a UI. We want to demonstrate the implementation of such kinds of UI with our concept by a simple desk calculator UI shown in Fig. 2. The implementation of this UI requires the access of the UI to some state that can be modified by the event handlers associated to the different buttons. In our application, the value of the state is a pair (d, f) containing the current operand d and an accumulator function f that is applied to d when the button "=" is pressed (this idea is due to [23]). In order to allow the change of the state's value by event handlers, we use `IORefs`, a concept from Haskell to deal with modifiable state. `IORefs` are references to stateful objects, where their states can only be accessed and changed by the predefined I/O actions `readIORef` and `writeIORef` (in order to ensure referential transparency). Thus, the calculator UI can be implemented as follows (where the parameter `stref` is an `IORef` to the calculator's state):

```
calcUI stref = col [entryS [Class [Bg Yellow]] display "0",
                   row (map cbutton ['1','2','3','+']),
                   row (map cbutton ['4','5','6','-']),
                   row (map cbutton ['7','8','9','*']),
                   row (map cbutton ['C','0','=','/'])]
where
  display free
  cbutton c = button (buttonPressed c) [c]
  buttonPressed c env = do
    state <- readIORef stref
    let (d,f) = processButton c state
    writeIORef stref (d,f)
    setValue display (show d) env
```

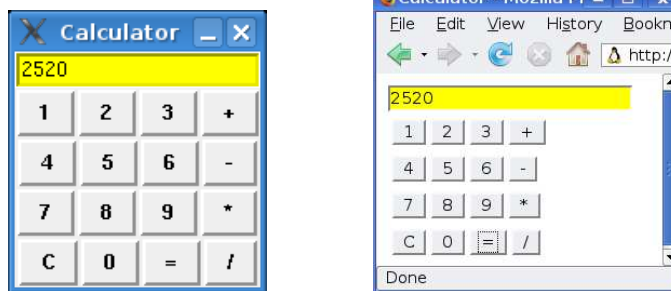


Fig. 2. A simple desk calculator UI executed as a GUI (left) and as a WUI (right)

The operator `entryS` is similar to `entry` but has a further argument to specify the initial layout of this widget (here: the background color). Note that we exploit the higher-order features of Curry to create the individual buttons by the generic function `cbutton` in a compact way. Each button has an associated event handler `buttonPressed` that reads the current state, modifies it, and shows the new operand in the entry widget referenced by the variable `display`. The actual update of the state depending on the selected button is computed by the operation `processButton`:

```
processButton :: Char -> (Int,Int->Int) -> (Int,Int->Int)
processButton b (d,f)
  | isDigit b = (10*d + ord b - ord '0', f)
  | b=='+'   = (0,((f d) +))
  | b=='-'   = (0,((f d) -))
  | b=='*'   = (0,((f d) *))
  | b=='/'   = (0,((f d) 'div'))
  | b=='='   = (f d, id)
  | b=='C'   = (0, id)
```

Finally, the complete application is executed by evaluating the operation `main` that first creates a new `IORef` object and then runs the UI with this object:

```
main = do stref <- newIORef (0,id)
         runUI "Calculator" (calcUI stref)
```

5 Implementation Issues

The definition of the components to specify a user interface, as discussed in Section 3, are contained in a library `UI` so that one has to import this library in order to define an interface. However, such an interface is not executable without specifying whether it should be run as a GUI or a WUI. For this purpose, our framework provides two implementations of the general UI concept by transforming UIs into GUIs or into WUIs. The necessary functionality is contained

in the libraries `UI2GUI` and `UI2HTML`, respectively. In order to execute a UI as a GUI (as shown in the left-hand sides of Fig. 1 and 2), one has to import the library `UI2GUI` (which has the same interface as `UI`) instead of `UI`, i.e., one has to put the import declaration

```
import UI2GUI
```

at the beginning of the module containing the corresponding UI specification. In order to execute a UI as a WUI (as shown in the right-hand sides of Fig. 1 and 2), one has to replace `UI2GUI` by `UI2HTML` in the import declaration, and everything else is left unchanged (apart from the command to generate an executable from the corresponding Curry program).

The implementation of the library `UI2GUI` is straightforward by exploiting the existing Curry library `GUI` [8] and mapping UI elements into corresponding GUI elements. The implementation of the library `UI2HTML` is more advanced since the existing Curry library `HTML` [9] does not support server interaction inside a web page. Since this is possible by the Ajax framework [6], we have added extensions to the `HTML` library (based on Ajax) to support the interaction model implied by the UI library.

6 Transforming GUIs into WUIs

Since the structure of UI elements is very similar to the elements of the Curry library `GUI`, which has been already used for various programming applications (e.g., [11, 16, 22]), one can also use our concept of UIs to enable the execution of such GUI-based applications as web applications. For this purpose, we have also implemented a library `GUI2HTML` that provides the same interface as the library `GUI` but executes a GUI as a WUI by exploiting the library `UI2HTML`. For instance, we have used this implementation to execute the Curry analysis environment `CurryBrowser` (its implementation consists of almost 4000 lines of Curry code), which is written in Curry and has a quite advanced graphical user interface (see [11]), in a standard web browser. The only necessary change was the replacement of the import of the library `GUI` by the import of the library `GUI2HTML` in the source code of the `CurryBrowser` implementation.

7 Conclusions

We described a framework to implement user interfaces in a high level, declarative manner. Our approach is based on separating the structural, functional, and layout aspects of a user interface. We showed that the features available in functional logic languages can be exploited to provide appropriate specifications of these issues. The hierarchical structure of UIs can be easily specified as term structures. The associated functionality can be specified by attaching event handlers (i.e., functions) to the elements of these term structures. The connections of event handlers to the individual widgets of the UI can be described by logic

variables. This avoids typical programming errors in untyped scripting languages and supports compositionality in the construction of complex UIs (e.g., there can be several instances of the same UI specification in a UI without name conflicts, see also [9]). Finally, the concrete layout is separated from the structural and functional aspects of the UI. This supports the use of the same UI specification in different contexts, i.e., one can create either graphical user interfaces on single local host computers or web-based user interfaces from such descriptions only by importing the appropriate libraries. This simplifies the programming efforts to combine the advantages of existing user interface technologies. Finally, our framework also enables the transformation of existing GUI applications into web applications and the embedding of UIs into arbitrary HTML pages (the latter aspect was not discussed in detail in this extended abstract).

The various features of the declarative base language Curry, in particular, algebraic data types, functions as first class citizens, logic variables, and polymorphic types, supports the high level, compact, and reliable specification of UIs (see also [8,9]) that can be used in different contexts. For future work it might be interesting to explore whether the same or a slightly modified concept can be also used to create user interfaces for other architectures, e.g., mobile devices.

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